

Semantic Networks and Description Logics

Description Logics – Terminology and Notation

Knowledge Representation and Reasoning

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Motivation

- ▶ Main problem with **semantic networks** and **frames**
- The lack of **formal semantics!**
- ▶ Disadvantage of simple **inheritance networks**
- Concepts are atomic and do not have any **structure**
- ↔ Brachman's **structural inheritance networks** (1977)

Description Logics – Terminology and Notation

Introduction

Concept and Roles

TBox and ABox

Reasoning Services

Outlook

Structural Inheritance Networks

- ▶ Concepts are **defined/described** using a small set of well-defined operators
- ▶ Distinction between **conceptual** and **object-related** knowledge
- ▶ Computation of **subconcept relation** and of **instance relation**
- ▶ **Strict inheritance** (of the entire structure of a concept)

Systems and Applications

► Systems:

- **KL-ONE**: First implementation of the ideas (1978)
- ... then **NIKL**, **KL-TWO**, **KRYPTON**, **KANDOR**, **CLASSIC**, **BACK**, **KRIS**, **YAK**, **CRACK** ...
- ... currently **FaCT**, **DLP**, **RACER** 1998

► Applications:

- First, natural language understanding systems
 - ... then configuration systems,
 - ... information systems,
 - ... currently, it is one tool for the *semantic web*
- ↪ **DAML+OIL**, now **OWL**

Description Logics

- Previously also *KL-ONE-like languages*, *frame-based languages*, *terminological logics*, *concept languages*
- **Description Logics (DL)** allow us
 - to describe concepts using *complex descriptions*,
 - to introduce the terminology of an application and to structure it (**TBox**),
 - to introduce objects (**ABox**) and relate them to the introduced terminology,
 - and to *reason* about the terminology and the objects.

Informal Example

Male is:	the opposite of female
A human is a kind of:	living entity
A woman is:	a human and a female
A man is:	a human and a male
A mother is:	a woman with at least one child that is a human
A father is:	a man with at least one child that is a human
A parent is:	a mother or a father
A grandmother is:	a woman, with at least one child that is a parent
A mother-wod is:	a mother with only male children

Elizabeth is a woman
 Elizabeth has the child Charles
 Charles is a Man
 Diana is a Mother-wod
 Diana has the Child William

Possible Questions:
 Is a grandmother a parent?
 Is Diana a parent?
 Is William a man?
 Is Elizabeth a mother-wod?

Atomic Concepts and Roles

- **Concept names:**
 - E.g., Grandmother, Male, ... (in the following usually *capitalized*)
 - We will use **symbols** such as A, A_1, \dots
 - **Semantics:** Monadic predicates $A(\cdot)$ or set-theoretically a subset of the universe $A^{\mathcal{I}} \subseteq \mathcal{D}$.
- **Role names:**
 - In our example, e.g., child. Often we will use names such as has-child or something similar (in the following usually *lowercase*).
 - Role names are *disjoint* from concept names
 - **Symbolically:** t, t_1, \dots
 - **Semantics:** Dyadic predicates $t(\cdot, \cdot)$ or set-theoretically $t^{\mathcal{I}} \subseteq \mathcal{D} \times \mathcal{D}$.

Concept and Role Description

- ▶ Out of *concept* and *role names*, complex **descriptions** can be created
- ▶ In our example, e.g. “a Human and Female.”
- ▶ **Symbolically**: C for concept descriptions and r for role descriptions
- ▶ Which particular constructs are available depends on the chosen description logic
- ▶ **Predicate logic semantics**: A concept descriptions C corresponds to a formula $C(x)$ with the free variable x . Similarly with r : It corresponds to formula $r(x, y)$ with free variables x, y .
- ▶ **Set semantics**:

$$C^{\mathcal{I}} = \{d \mid C(d) \text{ “is true in” } \mathcal{I}\}$$

$$r^{\mathcal{I}} = \{(d, e) \mid r(d, e) \text{ “is true in” } \mathcal{I}\}$$

Boolean Operators

- ▶ **Syntax**: let C and D be concept descriptions, then the following are also concept descriptions:
 - ▶ $C \sqcap D$ (**Concept conjunction**)
 - ▶ $C \sqcup D$ (**Concept disjunction**)
 - ▶ $\neg C$ (**Concept negation**)
- ▶ **Examples**:
 - ▶ Human \sqcap Female
 - ▶ Father \sqcup Mother
 - ▶ \neg Female
- ▶ **Predicate logic semantics**: $C(x) \wedge D(x)$, $C(x) \vee D(x)$, $\neg C(x)$
- ▶ **Set semantics**: $C^{\mathcal{I}} \cap D^{\mathcal{I}}$, $C^{\mathcal{I}} \cup D^{\mathcal{I}}$, $\mathcal{D} - C^{\mathcal{I}}$

Role Restrictions

- ▶ **Motivation**:
 - ▶ Often we want to describe something by *restricting* the possible “fillers” of a role, e.g. Mother-wod.
 - ▶ Sometimes we want to say that there is at least a filler of a particular type, e.g. Grandmother
- ▶ **Idea**: Use **quantifiers** that range over the role-fillers
 - ▶ Mother $\sqcap \forall$ has-child.Man
 - ▶ Woman $\sqcap \exists$ has-child.Parent
- ▶ **Predicate logic semantics**:

$$(\exists r.C)(x) = \exists y : (r(x, y) \wedge C(y))$$

$$(\forall r.C)(x) = \forall y : (r(x, y) \rightarrow C(y))$$

Set semantics:

$$(\exists r.C)^{\mathcal{I}} = \{d \mid \exists e : (d, e) \in r^{\mathcal{I}} \wedge e \in C^{\mathcal{I}}\}$$

$$(\forall r.C)^{\mathcal{I}} = \{d \mid \forall e : (d, e) \in r^{\mathcal{I}} \rightarrow e \in C^{\mathcal{I}}\}$$

Cardinality Restriction

- ▶ **Motivation**:
 - ▶ Often we want to describe something by *restricting the number* of possible “fillers” of a role, e.g., a Mother with at least 3 children or at most 2 children.
- ▶ **Idea**: We restrict the cardinality of the role filler sets:
 - ▶ Mother $\sqcap (\geq 3 \text{ has-child})$
 - ▶ Mother $\sqcap (\leq 2 \text{ has-child})$
- ▶ **Predicate logic semantics**:

$$(\geq n r)(x) = \exists y_1 \dots y_n : (r(x, y_1) \wedge \dots \wedge r(x, y_n) \wedge y_1 \neq y_2 \wedge \dots \wedge y_{n-1} \neq y_n)$$

$$(\leq n r)(x) = \neg(\geq n + 1 r)(x)$$

- ▶ **Set semantics**:

$$(\geq n r)^{\mathcal{I}} = \{d \mid |\{e \mid r^{\mathcal{I}}(d, e)\}| \geq n\}$$

$$(\leq n r)^{\mathcal{I}} = \mathcal{D} - (\geq n + 1 r)^{\mathcal{I}}$$

Inverse Roles

► **Motivation:**

- How can we describe the concept “*children of rich parents*”?

► **Idea:** Define the “inverse” role for a given role (the **converse relation**)

- has-child^{-1}

► **Application:** $\exists \text{has-child}^{-1}.\text{Rich}$

► **Predicate logic semantics:**

$$r^{-1}(x, y) = r(y, x)$$

► **Set semantics:**

$$(r^{-1})^{\mathcal{I}} = \{(d, e) \mid (e, d) \in r^{\mathcal{I}}\}$$

Role Composition

► **Motivation:**

- How can we define the role `has-grandchild` given the role `has-child`?

► **Idea:** Compose roles (as one can compose binary relations)

- $\text{has-child} \circ \text{has-child}$

► **Predicate logic semantics:**

$$(r \circ s)(x, y) = \exists z : (r(x, z) \wedge s(z, y))$$

► **Set semantics:**

$$(r \circ s)^{\mathcal{I}} = \{(d, e) \mid \exists f : (d, f) \in r^{\mathcal{I}} \wedge (f, e) \in s^{\mathcal{I}}\}$$

Role Value Maps

► **Motivation:**

- How do we express the concept “*women, who know all the friends of their children*”

► **Idea:** Relate role filler sets to each other

- $\text{Woman} \sqcap (\text{has-child} \circ \text{has-friend} \sqsubseteq \text{knows})$

► **Predicate logic semantics:**

$$(r \sqsubseteq s)(x) = \forall y : (r(x, y) \rightarrow s(x, y))$$

► **Set semantics:** Let $r^{\mathcal{I}}(d) = \{e \mid r^{\mathcal{I}}(d, e)\}$.

$$(r \sqsubseteq s)^{\mathcal{I}} = \{d \mid r^{\mathcal{I}}(d) \subseteq s^{\mathcal{I}}(d)\}$$

► **Note:** Role value maps lead to undecidability of satisfiability of concept descriptions!

Terminology Box

► In order to **introduce** new terms, we use two kinds of **terminological axioms**:

- $A \doteq C$
- $A \sqsubseteq C$

where A is a **concept name** and C is a **concept description**.

► A **terminology** or **TBox** is a finite set of such axioms with the following additional restrictions:

- no multiple definitions of the same symbol such as $A \doteq C$, $A \sqsubseteq D$
- no cyclic definitions (even not indirectly), such as $A \doteq \forall r.B$, $B \doteq \exists s.A$

TBoxes: Semantics

- ▶ TBoxes restrict the set of possible interpretations.
- ▶ **Predicate logic semantics:**
 - ▶ $A \doteq C$ corresponds to $\forall x : (A(x) \leftrightarrow C(x))$
 - ▶ $A \sqsubseteq C$ corresponds to $\forall x : (A(x) \rightarrow C(x))$
- ▶ **Set semantics:**
 - ▶ $A \doteq C$ corresponds to $A^{\mathcal{I}} = C^{\mathcal{I}}$
 - ▶ $A \sqsubseteq C$ corresponds to $A^{\mathcal{I}} \subseteq C^{\mathcal{I}}$
- ▶ Non-empty interpretations which satisfy all terminological axioms are called **models** of the TBox.

Assertional Box

- ▶ In order to state something about objects in the world, we use two forms of **assertions**:
 - ▶ $a : C$
 - ▶ $(a, b) : r$
 where a and b are **individual names** (e.g., ELIZABETH, PHILIP), C is a **concept description**, and r is a **role description**.
- ▶ An **ABox** is a finite set of assertions.

ABoxes: Semantics

- ▶ **Individual names** are interpreted as elements of the universe under the **unique-name-assumption**, i.e., different names refer to different objects.
- ▶ **Assertions** express that an object is an instance of a concept or that two objects are related by a role.
- ▶ **Predicate logic semantics:**
 - ▶ $a : C$ corresponds to $C(a)$
 - ▶ $(a, b) : r$ corresponds to $r(a, b)$
- ▶ **Set semantics:**
 - ▶ $a^{\mathcal{I}} \in D$
 - ▶ $a : C$ corresponds to $a^{\mathcal{I}} \in C^{\mathcal{I}}$
 - ▶ $(a, b) : r$ corresponds to $(a^{\mathcal{I}}, b^{\mathcal{I}}) \in r^{\mathcal{I}}$
- ▶ **Models** of an ABox and of ABox+TBox can be defined analogously to models of a TBox.

Example TBox

```

Male  $\doteq$   $\neg$ Female
Human  $\sqsubseteq$  Living_entity
Woman  $\doteq$  Human  $\sqcap$  Female
Man  $\doteq$  Human  $\sqcap$  Male
Mother  $\doteq$  Woman  $\sqcap$   $\exists$ has-child.Human
Father  $\doteq$  Man  $\sqcap$   $\exists$ has-child.Human
Parent  $\doteq$  Father  $\sqcup$  Mother
Grandmother  $\doteq$  Woman  $\sqcap$   $\exists$ has-child.Parent
Mother-without-daughter  $\doteq$  Mother  $\sqcap$   $\forall$ has-child.Male
Mother-with-many-children  $\doteq$  Mother  $\sqcap$  ( $\geq 3$ has-child)

```

Example ABox

```

CHARLES: Man
EDWARD: Man
ANDREW: Man
DIANA: Mother-without-daughter
(ELIZABETH, CHARLES): has-child
(ELIZABETH, EDWARD): has-child
(ELIZABETH, ANDREW): has-child
(DIANA, WILLIAM): has-child
(CHARLES, WILLIAM): has-child

DIANA: Woman
ELIZABETH: Woman

```





Outlook

- ▶ Can we **reduce** the reasoning services to perhaps just one problem?
- ▶ What could be **reasoning algorithms**?
- ▶ What about **complexity** and **decidability**?
- ▶ What has all that to do with **modal logics**?
- ▶ How can one build **efficient systems**?

Some Reasoning Services

- ▶ Does a description C make sense at all, i.e., is it **satisfiable**?
- ↔ A concept description C is satisfiable iff there exists an interpretation \mathcal{I} such that $C^{\mathcal{I}} \neq \emptyset$.
- ▶ Is one concept a specialization of another one, is it **subsumed**?
- ↔ C is **subsumed by** D , in symbols $C \sqsubseteq D$ iff we have for all interpretations $C^{\mathcal{I}} \subseteq D^{\mathcal{I}}$.
- ▶ Is a an **instance** of a concept C ?
- ↔ a is an instance of C iff for all interpretations, we have $a^{\mathcal{I}} \in C^{\mathcal{I}}$.
- **Note:** These questions can be posed with or without a TBox that restricts the possible interpretations.

Literature

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Summary: Concept Descriptions

Abstract	Concrete	Interpretation
A	A	$A^{\mathcal{I}}$
$C \sqcap D$	(and $C D$)	$C^{\mathcal{I}} \cap D^{\mathcal{I}}$
$C \sqcup D$	(or $C D$)	$C^{\mathcal{I}} \cup D^{\mathcal{I}}$
$\neg C$	(not C)	$\mathcal{D} - C^{\mathcal{I}}$
$\forall r.C$	(all $r C$)	$\{d \in \mathcal{D} \mid r^{\mathcal{I}}(d) \subseteq C^{\mathcal{I}}\}$
$\exists r$	(some r)	$\{d \in \mathcal{D} \mid r^{\mathcal{I}}(d) \neq \emptyset\}$
$\geq n r$	(atleast $n r$)	$\{d \in \mathcal{D} \mid r^{\mathcal{I}}(d) \geq n\}$
$\leq n r$	(atmost $n r$)	$\{d \in \mathcal{D} \mid r^{\mathcal{I}}(d) \leq n\}$
$\exists r.C$	(some $r C$)	$\{d \in \mathcal{D} \mid r^{\mathcal{I}}(d) \cap C^{\mathcal{I}} \neq \emptyset\}$
$\geq n r.C$	(atleast $n r C$)	$\{d \in \mathcal{D} \mid r^{\mathcal{I}}(d) \cap C^{\mathcal{I}} \geq n\}$
$\leq n r.C$	(atmost $n r C$)	$\{d \in \mathcal{D} \mid r^{\mathcal{I}}(d) \cap C^{\mathcal{I}} \leq n\}$
$r = s$	(eq $r s$)	$\{d \in \mathcal{D} \mid r^{\mathcal{I}}(d) = s^{\mathcal{I}}(d)\}$
$r \neq s$	(neq $r s$)	$\{d \in \mathcal{D} \mid r^{\mathcal{I}}(d) \neq s^{\mathcal{I}}(d)\}$
$r \sqsubseteq s$	(subset $r s$)	$\{d \in \mathcal{D} \mid r^{\mathcal{I}}(d) \subseteq s^{\mathcal{I}}(d)\}$
$g = h$	(eq $g h$)	$\{d \in \mathcal{D} \mid g^{\mathcal{I}}(d) = h^{\mathcal{I}}(d) \neq \emptyset\}$
$g \neq h$	(neq $g h$)	$\{d \in \mathcal{D} \mid \emptyset \neq g^{\mathcal{I}}(d) \neq h^{\mathcal{I}}(d) \neq \emptyset\}$
$\{i_1, i_2, \dots, i_n\}$	(oneof $i_1 \dots i_n$)	$\{i_1^{\mathcal{I}}, i_2^{\mathcal{I}}, \dots, i_n^{\mathcal{I}}\}$

Summary: Role Descriptions

Abstract	Concrete	Interpretation
t	t	$t^{\mathcal{I}}$
f	f	$f^{\mathcal{I}}$, (functional role)
$r \sqcap s$	(and $r s$)	$r^{\mathcal{I}} \cap s^{\mathcal{I}}$
$r \sqcup s$	(or $r s$)	$r^{\mathcal{I}} \cup s^{\mathcal{I}}$
$\neg r$	(not r)	$\mathcal{D} \times \mathcal{D} - r^{\mathcal{I}}$
r^{-1}	(inverse r)	$\{(d, d') \mid (d', d) \in r^{\mathcal{I}}\}$
$r \upharpoonright_C$	(restr $r C$)	$\{(d, d') \in r^{\mathcal{I}} \mid d' \in C^{\mathcal{I}}\}$
r^+	(trans r)	$(r^{\mathcal{I}})^+$
$r \circ s$	(compose $r s$)	$r^{\mathcal{I}} \circ s^{\mathcal{I}}$
$\mathbf{1}$	self	$\{(d, d) \mid d \in \mathcal{D}\}$